Proof Reconstruction: Translating Proofs

Alejandro Gómez-Londoño Advisor - Andrés Sicard-Ramírez

EAFIT University

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Introduction A (very) general idea of the context





An interactive prover is a software tool aiding the development of formal proofs by man-machine collaboration.¹

¹Matita development team, Matita website, http://matita.cs.unibo.it/index.shtml

Deals with the development of computer programs that show that some statement (the conjecture) is a logical consequence of a set of statements (the axioms and hypotheses).²

 $^{^{2}} http://www.cs.miami.edu/{\sim}tptp/OverviewOfATP.html$

Introduction ATPs input/output



³Sutcliffe, G. The TPTP Problem Library and Associated In-frastructure: The FOF and CNF Parts. 2009.

ATPs:

- Vampire
- E E
- Metis
- SPASS
- Equinox

Proof assistants:

- Coq
- Agda
- Isabelle
- Mizar
- NuPRL

Introduction



Introduction



Introduction



⁴Jasmin C. Blanchette, Cezary Kaliszyk and Lawrence C. Paulson, Hammering towards QED. 2014.

Hand written proof

$$\frac{x \quad y}{x \land y} \quad x \land y \Rightarrow z$$

TPTP problem

fof(a_0,axiom,x).
fof(a_1,axiom,y).
fof(a_2,axiom, ((x & y) => z)).
fof(c_0,conjecture, z).

Proof reconstruction Example

TSTP proof

```
fof(s_1,plain,(z),
    inference(modus_ponens,[],[a_2,s_0])).
```

Proof reconstruction Example

■ Agda proof --conjunction data _^_ (P : Set) (Q : Set) : Set where ^-intro : P ^ Q ^ (P ^ Q) proof : { X Y Z : Set} → X → Y → (X ^ Y → Z) → Z proof x y f = f (^-intro x y)

Proof reconstruction



```
fof(a 0,axiom,x).
fof(a 1,axiom,y).
fof(a 2, axiom, ((x \& y) \Rightarrow z)).
fof(c 0, conjecture, z).
fof(s 0, plain, (x \& y),
    inference(conjunction,[],[a 0,a 1])).
fof(s 1,plain,(z),
    inference(modus ponens,[],[a_2,s_0])).
fof(r 0,plain,($true),
    inference(simplify,[],[s 1,c 0])).
```

Proof reconstruction Parser and AST construction

```
F {name = "s 0",
   role = Plain,
   formula = "x" (:\&:) "y",
  annotations = Conjunction ["a 0", "a 1"]
  },
F {name = "s 1",
   role = Plain,
   formula = "z",
   annotations = ModusPonens ["a 2", "s 0"]
 }.
F \{name = "r 0",
   role = Plain,
   formula = "$True",
  annotations = Simplify ["s 1", "c 0"]
  }
```

Proof reconstruction DAG



Haskell and Agda were chosen as the programing languages for the implementation.

- Haskell was used for parsing and AST construction
- In Agda we will create and analyze the DAG.

Past results

Metis⁵ was chosen as our ATP

- Uses TPTP as input format.
- Outputs proofs in TSTP format.
- Each refutation step is one of 6 rules.
- Has respectable performance.

⁵Joe Hurd. First-Order Proof Tactics in Higher-Order Logic Theorem Provers. 2003.

A modified version of the logic-tptp⁶ Haskell library was used to implement a TSTP parser capable of analyze Metis proofs.

This project is freely available on github⁷.

⁶https://hackage.haskell.org/package/logic-TPTP
⁷https://github.com/agomezl/tstp2agda

Translate into idiomatic Agda code the AST resulting from the parsing of an ATP-generated proof.

- Translate to Agda code the Haskell AST data type.
- Build an Agda library that implements the logical kernel of the ATP.
- Reconstruct the proof in Agda using the aforementioned library.