

# CM0845 Logic

## Propositional Logic: Satisfiability

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# Propositional Logic: Satisfiability

## Remark

The reference for this section is Ben-Ari [2012, § 2.5].

# Satisfiability, Validity, Unsatisfiability and Falsifiability

Let  $\varphi \in \text{PROP}$ .

## Definitions

- (i)  $\varphi$  is **satisfiable** iff  $\llbracket \varphi \rrbracket_v = 1$  for some interpretation  $v$ .  
In this case,  $v$  is called a model for  $\varphi$ .
- (ii)  $\varphi$  is **valid** (a tautology), denoted  $\models \varphi$ , iff  $\llbracket \varphi \rrbracket_v = 1$  for all interpretations  $v$ .
- (iii)  $\varphi$  is **unsatisfiable** iff it is not satisfiable, that is, if  $\llbracket \varphi \rrbracket_v = 0$  for all interpretations  $v$ .
- (iv)  $\varphi$  is **falsifiable**, denoted  $\not\models$ , iff it is not valid, that is, if  $\llbracket \varphi \rrbracket_v = 0$  for some interpretation  $v$ .

# Satisfiability, Validity, Unsatisfiability and Falsifiability

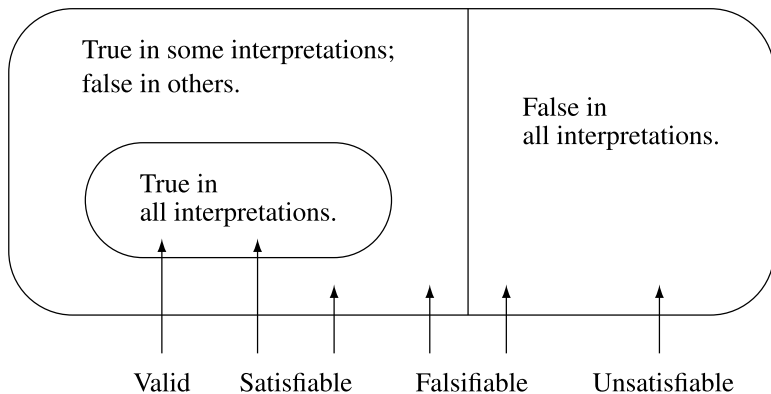


Figure 2.6 of [Ben-Ari 2012].

# Satisfiability, Validity, Unsatisfiability and Falsifiability

Theorem (Ben-Ari [2012], Theorem 2.39)

Let  $\varphi \in \text{PROP}$ .

- (i) The proposition  $\varphi$  is valid if and only if  $\neg\varphi$  is unsatisfiable.
- (ii) The proposition  $\varphi$  is satisfiable if and only if  $\neg\varphi$  is falsifiable.

# Satisfiability of a Set of Propositions

Let  $\Gamma = \{\varphi_1, \dots\}$  be a set of propositions.

## Definitions

- (i)  $\Gamma$  is **satisfiable** iff there exists an interpretation  $v$  such that  $\llbracket \varphi \rrbracket_v = 1$  for all  $\varphi_i \in \Gamma$ . In this case,  $v$  is a model of  $\Gamma$ .
- (ii)  $\Gamma$  is **unsatisfiable** iff for every interpretation  $v$ , there exists an  $\varphi_i \in \Gamma$  such that  $\llbracket \varphi \rrbracket_v = 0$ .

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- (ii)  $\Gamma$  is **unsatisfiable** iff for every interpretation  $v$ , there exists an  $\varphi_i \in \Gamma$  such that  $\llbracket \varphi \rrbracket_v = 0$ .

## Example

Prove that if  $\Gamma$  is unsatisfiable and for some  $i$ , the proposition  $\varphi_i$  is valid, then  $\Gamma - \{\varphi_i\}$  is unsatisfiable [Ben-Ari 2012, Exercise 2.15, p. 46].

# References



Ben-Ari, Mordechai [1993] (2012). *Mathematical Logic for Computer Science*. 3rd ed. Springer (cit. on pp. 2, 4–7).