# CM0246 Discrete Structures Euler and Hamilton Paths 

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## Preliminaries

## Convention

The number assigned to chapters, examples, exercises, figures, sections, and theorems on these slides correspond to the numbers assigned in the textbook (Rosen 2004).

## The Problem of the Seven Bridges of Königsberg

## Problem

It is possible to start at some location in the town, travel across all the bridges once without crossing any bridge twice, and return to the starting point? ${ }^{\dagger}$

${ }^{\dagger}$ Figure source: (Rosen 2012, § 10.5, Fig. 1).

## Paths

## Definition

Let $n$ be a non-negative integer and $G$ an undirected graph. A path of length $n$ from $u$ to $v$ in $G$ is a sequence of $n$ edges $e_{1}, e_{2}, \ldots, e_{n}$ of $G$ such that $f\left(e_{1}\right)=\left\{x_{0}, x_{1}\right\}, f\left(e_{2}\right)=\left\{x_{1}, x_{2}\right\}, \ldots, f\left(e_{n}\right)=\left\{x_{n-1}, x_{n}\right\}$, where $x_{0}=u$ and $x_{n}=v$.

Example
Whiteboard.

## Paths

## Definition

A path is a circuit if $u=v$ (it begins and ends at the same vertex) and has length greater than zero.

Example
Whiteboard.

## Paths

## Definition

A path is a circuit if $u=v$ (it begins and ends at the same vertex) and has length greater than zero.

## Example

Whiteboard.
Definition
A path/circuit is simple if it does not contain the same edge more than once.

Example
Whiteboard.

## Connectedness

## Definition

An undirected graph is called connected (conexo) if there is a path between every pair of distinct vertices of the graph.

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Example

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## Euler Paths

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## Exercise

Find an Euler path in the following graph:


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## Exercise

Find an Euler path in the following graph:


Solution
An Euler path is $a, c, d, e, b, d, a, b$.

## Euler Circuits

## Definition

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An Euler circuit in a graph $G$ is a simple circuit containing every edge of $G$.

## Exercise

Find an Euler circuit in the following graph:


## Solution

An Euler circuit is $a, e, c, d, e, b, a$.

## Euler Paths and Circuits

Theorem 1 (p. 543)
A connected multigraph with at least two vertices has an Euler circuit, if and only if, each of its vertices has even degree.

## Euler Paths and Circuits

Example (solution to the problem of the seven bridges of Königsberg)
It is possible to start at some location in the town, travel across all the bridges once without crossing any bridge twice, and return to the starting point?

## Euler Paths and Circuits

Example (solution to the problem of the seven bridges of Königsberg)
It is possible to start at some location in the town, travel across all the bridges once without crossing any bridge twice, and return to the starting point?

## Solution

No, it is not possible, because the multigraph representing the bridges has (four) vertices of odd degree. ${ }^{\dagger}$

${ }^{\dagger}$ Figure source: (Rosen 2012, § 10.5, Fig. 1).

## Euler Paths and Circuits

Theorem 2 (p. 544)
A connected multigraph has an Euler path but not an Euler circuit, if and only if it has exactly two vertices of odd degree.

## Euler Paths and Circuits

## Exercise

Has the following graph an Euler path? If so, find it.


## Euler Paths and Circuits

## Exercise

Has the following graph an Euler path? If so, find it.


## Solution

The only vertices with degree odd are $b$ and $d$, so the graph has an Euler path. An Euler path is

$$
b, a, g, f, e, d, c, g, b, c, f, d
$$

## Euler Paths and Circuits

Problem 4 (p. 550)
Determine whether the given graph has an Euler circuit. Construct such a circuit when one exists. If no Euler circuit exists, determine whether the graph has an Euler path and construct such a path if one exists.


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Determine whether the given graph has an Euler circuit. Construct such a circuit when one exists. If no Euler circuit exists, determine whether the graph has an Euler path and construct such a path if one exists.

- The graph has no an Euler
 circuit because it has at least a vertice of odd degree $(\delta(f)=3)$.
- The graph has an Euler path because the only two vertices of odd degree are $f$ and $c$.
- Euler path:
$f, a, e, f, b, a, d, e, c, d, b, c$.


## Euler Paths and Circuits

## Problem 9 (p. 550)

Suppose that in addition to the seven bridges of Königsberg there were two additional bridges, connecting regions $B$ and $C$ and regions $B$ and $D$, respectively. Could someone cross all nine of these bridges exactly once and return to the starting point? ${ }^{\dagger}$


[^0]
## Euler Paths and Circuits

## Problem 9 (p. 550)

Suppose that in addition to the seven bridges of Königsberg there were two additional bridges, connecting regions $B$ and $C$ and regions $B$ and $D$, respectively. Could someone cross all nine of these bridges exactly once and return to the starting point? ${ }^{\dagger}$


The problem asks for building an Euler circuit. It it not possible because the graph representing the problem has vertices of odd degree (e.g. $\delta(A)=5$ ).

[^1]
## The Icosian Puzzle

The icosian puzzle
‘The icosian puzzle (invented by William Rowan Hamilton) consisted of a wooden dodecahedron (a polyhedron with 12 regular pentagons as faces), with a peg at each vertex of the dodecahedron, and string. The 20 vertices of the dodecahedron were labeled with different cities in the world. The object of the puzzle was to start at a city and travel along the edges of the dodecahedron, visiting each of the other 19 cities exactly once, and end back at the first city.' (Rosen 2012, 7th ed. § 10.5, pp. 698-699).


## Hamilton Paths

## Definition

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## Exercise

Find a Hamilton path in the following graph:


Solution
A Hamilton path is $a, b, c, d$.

## Hamilton Circuits

## Definition

A simple circuit in a graph $G$ that passes through every vertex exactly once is called a Hamilton circuit.

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## Exercise

Find a Hamilton circuit in the following graph:


## Hamilton Circuits

## Definition

A simple circuit in a graph $G$ that passes through every vertex exactly once is called a Hamilton circuit.

## Exercise

Find a Hamilton circuit in the following graph:


Solution
A Hamilton circuit is $a, b, c, d, e, a$.

## Hamilton Circuits

Example
Solution to the icosian puzzle. ${ }^{\dagger}$

${ }^{\dagger}$ Figure source: (Rosen 2012, § 10.5, Fig. 9).

## Hamilton Circuits

Theorem 3 (Dirac's Theorem, p. 548)
If $G$ is a simple graph with $n$ vertices with $n \geq 3$ such that the degree of every vertex in $G$ is at least $n / 2$, then $G$ has a Hamilton circuit.

## Hamilton Circuits

Theorem 3 (Dirac's Theorem, p. 548)
If $G$ is a simple graph with $n$ vertices with $n \geq 3$ such that the degree of every vertex in $G$ is at least $n / 2$, then $G$ has a Hamilton circuit.

Theorem 4 (Ore's Theorem, p. 548)
If $G$ is a simple graph with $n$ vertices with $n \geq 3$ such that $\delta(u)+\delta(v) \geq n$ for every pair of non-adjacent vertices $u$ and $v$ in $G$, then $G$ has a Hamilton circuit.

## References

Rosen, K. H. (2004). Matemática Discreta y sus Aplicaciones. 5th ed. Translated by José Manuel Pérez Morales and others. McGraw-Hill (cit. on p. 2).

- (2012). Discrete Mathematics and Its Applications. 7th ed. McGrawHill (cit. on pp. 3, 16, 17, 23-25, 32).


[^0]:    ${ }^{\dagger}$ Figure source: (Rosen 2012, § 10.5, Fig. 1).

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